

Stuxnet drivers: detailed analysis

artemonsecurity.blogspot.de/2017/04/stuxnet-drivers-detailed-analysis.html

There has passed already a lot of time since the publication of various detailed researches about Stuxnet and its components. All top AV vendors wrote own comprehensive papers, which reveal major information about destructive Stuxnet features. Some information about Stuxnet rootkits were published by Kaspersky [here](#), Symantec [here](#), ESET [here](#). However, the published information is not complete, because each of these documents covers only a specific sample of the rootkit and describes some of its functions. For example, Kaspersky analysis tries to summarize information about known Stuxnet drivers, but it doesn't contain any technical info about it. Another mentioned report from ESET contains information about two Stuxnet drivers, but this is not sufficient for complete summarizing.



First of all, it is need to be clear that from point of view of undocumented Windows kernel exploration, there are no something really interesting in Stuxnet drivers. I mean nothing interesting comparing with such advanced & sophisticated "civilian" rootkits like ZeroAccess or TDL4. These instances can be deeply embedded into a system, bypassing anti-rootkits and deceive low-level disk access tools. In contrast to them, authors of Stuxnet rootkits do not use such deep persistence into a compromised system. This analysis tries to summarize technical information about Stuxnet drivers.

As a starting point of our research, we can take already published information about Stuxnet drivers by Kaspersky. Their analysis Stuxnet/Duqu: The Evolution of Drivers summarizes some information about drivers that have been used by Stuxnet authors in cyber attacks.

Driver 1

File name: MRxCIs.sys

SHA256: 817a7f28a0787509c2973ce9ae85a95beb979e30b7b08e64c66d88372aa3da86

File size: 19840 bytes

Signed: No

Timestamp: 2009-01-01 18:53:25

Device object name: \Device\MRxClsDvX

Main purpose: code injection

AV detection ratio: 53/61

First driver contains sensitive text information such as rootkit device name and path to its service into registry as encrypted data. After starting, the driver performs decryption of this data and we can extract it. Note that name of rootkit service is almost matches its device object name. First dword of decrypted data is also interesting, because it stores some flags, which have an impact on driver behaviour. For example, first bit of this dword restricts the work of rootkit code into Windows safe mode, while second is used as anti-debug trick. If second bit and *ntoskrnl!KdDebuggerEnabled* are active, the driver will not load.

Decrypted rootkit data also stores name of registry value (Data) that is used by the rootkit to determining what files should be injected into processes. So, these decrypted data are stored in the next sequence.

\REGISTRY\MACHINE\SYSTEM\CurrentControlSet\Services\MRxCls

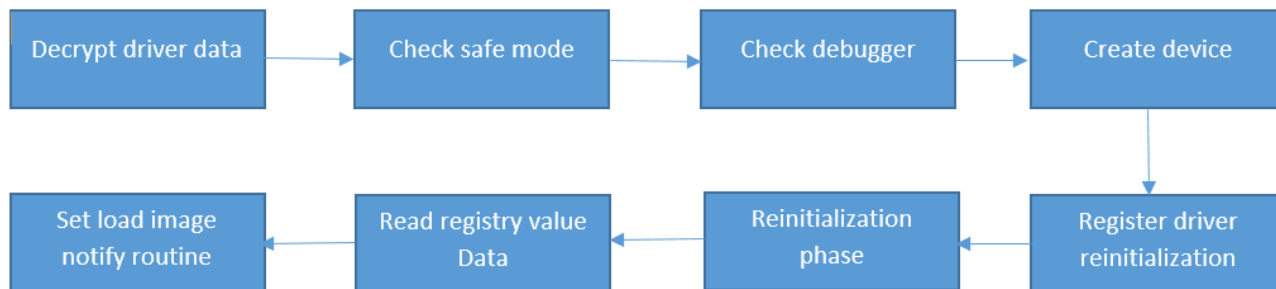
Data

\Device\MRxClsDvX

```
.text:0001041E BE 78 02 00 00      mov     esi, 278h
.text:00010423 B9 99 3E 01 00      mov     ecx, offset dwFlag
.text:00010428 E8 15 18 00 00      call   fnDecryptData ; ecx->data; esi->size
.text:0001042D 88 1D 98 3E 01+    mov     byte_13E98, bl
.text:00010433
.text:00010433          jCheckFlag: ; CODE XREF: start+70fj
.text:00010433 A1 99 3E 01 00      mov     eax, dwFlag
.text:00010438 A8 01              test    al, 1
.text:0001043A 74 10              jz     short jCheckRemoteDebugger
.text:0001043C A1 B0 23 01 00      mov     eax, ds:InitSafeBootMode
.text:00010441 39 18              cmp     [eax], ebx
.text:00010443 74 07              jz     short jCheckRemoteDebugger
.text:00010445
.text:00010445          jRetWithError: ; CODE XREF: start+B2lj
.text:00010445 B8 01 00 00 C0      mov     eax, 0C0000001h
.text:0001044A EB 14              jmp    short loc_10460
.text:0001044C          ; -----
.text:0001044C          jCheckRemoteDebugger: ; CODE XREF: start+90fj
.text:0001044C          ; start+99fj
.text:0001044C A1 99 3E 01 00      mov     eax, dwFlag
.text:00010451 A8 02              test    al, 2
.text:00010453 74 09              jz     short loc_1045E
.text:00010455 A1 AC 23 01 00      mov     eax, ds:KdDebuggerEnabled
.text:0001045A 38 18              cmp     [eax], bl
.text:0001045C 75 E7              jnz    short jRetWithError
.text:0001045F
```

As driver is registered by Stuxnet with "boot" loading type, it can't perform whole initialization in *DriverEntry*, because neither NT kernel nor file system is ready to perform requests from clients. So, it calls API *IoRegisterDriverReinitialization* and delays own initialization.

After *ReInitialize* rootkit function gets control, it checks Windows NT version and fills some dynamic imports. It also doing some preparatory operations and calls *PsSetLoadImageNotifyRoutine* for registering own handler on load image. This handler will response for code injection into processes. Below you can see scheme of rootkit initialization.



The driver registers following IRP handlers.

- IRP_MJ_CREATE
- IRP_MJ_CLOSE
- IRP_MJ_DEVICE_CONTROL (*fnDispatchIrpMjDeviceControl*)

If you are not familiar with Windows NT drivers development, it is worth to note that any driver that allows to open handles on its device, registers, at least, two IRP handlers: IRP_MJ_CREATE for supporting operations *ZwCreateFile* and IRP_MJ_CLOSE for *ZwClose*. Our driver supports *ZwDeviceIoControl* interface, that's why it registers IRP_MJ_DEVICE_CONTROL handler.

```

.text:00010A04          fnDispatchIrpMjDeviceControl proc near ; DATA XREF: DriverEntry+F310
.text:00010A04
.text:00010A04          var_1C                = dword ptr -1Ch
.text:00010A04          var_4                 = dword ptr -4
.text:00010A04          IRP                   = dword ptr 0Ch
.text:00010A04          esi_IRP = esi
.text:00010A04          6A 0C                push    0Ch
.text:00010A06          68 E8 3D 01 00       push    offset unk_13DE8
.text:00010A0B          E8 F0 17 00 00       call   fnSEH
.text:00010A0B
.text:00010A10          89 02 00 00 C0       mov     ecx, 0C0000002h
.text:00010A15          89 4D E4             mov     [ebp+var_1C], ecx
.text:00010A18          83 65 FC 00         and     [ebp+var_4], 0
.text:00010A1C          8B 75 0C            mov     esi_IRP, [ebp+IRP]
.text:00010A1F          8B 46 60            mov     eax, [esi_IRP+60h] ; IRP->IoStack
.text:00010A22          81 78 0C 00 38+     cmp     dword ptr [eax+0Ch], 223800h ; IoStack->DeviceIoControl.IoControlCode
.text:00010A29          74 04              jz     short jDispatchIOCTL
.text:00010A29
.text:00010A2B          8B C1              mov     eax, ecx
.text:00010A2D          EB 06              jmp     short loc_10A35
.text:00010A2D
.text:00010A2F          ; -----
.text:00010A2F          jDispatchIOCTL:    ; CODE XREF: fnDispatchIrpMjDeviceControl+251j
.text:00010A2F          56                push    esi_IRP
.text:00010A30          E8 13 10 00 00       call   fnDispatchIOCTL
.text:00010A30
.text:00010A30
.text:00010A35          loc_10A35:        ; CODE XREF: fnDispatchIrpMjDeviceControl+291j
.text:00010A35          8B F8              mov     edi, eax
.text:00010A37          89 7D E4             mov     [ebp-1Ch], edi
.text:00010A3A          83 4D FC FF         or     dword ptr [ebp-4], 0FFFFFFFh
.text:00010A3E          EB 11              jmp     short jCheckOnStatusPending ; STATUS_PENDING

```

handler supports pending I/O operation

Handler *fnDispatchIrpMjDeviceControl* serves only for one purpose: to call undocumented Windows NT function *ZwProtectVirtualMemory*. Client should send to driver special IOCTL code 0x223800 for that (*DeviceIoControl*) and provide a special prearranged structure with parameters for API call. The driver uses buffered I/O.

```

.text:00011A70
.text:00011A70
.text:00011A70 83 65 FC 00      jGetZwProtectVirtualMemoryAddr:      ; CODE XREF: fnDispatchIOCTL+2C1j
.text:00011A81 56              and     [ebp+var_4], 0
.text:00011A82 57              push   esi
.text:00011A83 57              push   esi
.text:00011A83 80 75 FC        lea    esi, [ebp+var_4]
.text:00011A86 80 7D F8        lea    edi, [ebp+pZwProtectVirtualMemory]
.text:00011A89 E8 94 EE FF FF  call   fnGetZwProtectVirtualMemoryAddr
.text:00011A89
.text:00011A8E 88 45 FC        mov     eax, [ebp+var_4]
.text:00011A91 85 C0          test   eax, eax
.text:00011A93 5F              pop     edi
.text:00011A94 5E              pop     esi
.text:00011A95 75 26          jnz    short jRet_
.text:00011A95
.text:00011A97 8D 43 20        lea    eax, [ebx+20h]
.text:00011A9A 50              push   eax ; 01dProtect
.text:00011A9B FF 30          push   dword ptr [eax] ; NewProtectWin32
.text:00011A9D 8D 43 18        lea    eax, [ebx+18h]
.text:00011AA0 50              push   eax ; RegionSize
.text:00011AA1 8D 43 10        lea    eax, [ebx+10h]
.text:00011AA4 50              push   eax ; BaseAddress
.text:00011AA5 FF 73 08        push   dword ptr [ebx+8] ; ProcessHandle
.text:00011AA8 88 45 F8        mov     eax, [ebp+pZwProtectVirtualMemory]
.text:00011AAB FF 10          call   dword ptr [eax] ; ZwProtectVirtualMemory
.text:00011AAD 85 C0          test   eax, eax
.text:00011AAF 75 0C          jnz    short jRet_

```

```

^
NTSTATUS
NTProtectVirtualMemory(
    _In HANDLE ProcessHandle,
    _Inout PVOID *BaseAddress,
    _Inout PSIZE_T RegionSize,
    _In WIN32_PROTECTION_MASK NewProtectWin32,
    _Out PULONG OldProtect
)
/****
Routine Description:

This routine changes the protection on a region of committed pages
within the virtual address space of the subject process. Setting
the protection on a range of pages causes the old protection to be
replaced by the specified protection value.

Note if a virtual address is locked in the working set and the
protection is changed to no access, the page is removed from the
working set since valid pages can't be no access.

Arguments:
    ProcessHandle - An open handle to a process object.

```

As we know, function *ZwProtectVirtualMemory* is not exported by the Windows kernel and this is another task which authors of MRXCLS.sys have been solved. For example, in case of Windows 2000, they try to find function signature with analysis of executable sections of ntoskrnl image. This signature you can see below.

```

.rdata:00012450
.rdata:00012450
.rdata:00012450
.rdata:00012450
.rdata:00012450 B8 77 00 00 00      mov     eax, 77h
.rdata:00012455 8D 54 24 04      lea    edx, [esp+4]
.rdata:00012459 CD 2E            int     2Eh ; DOS 2+ internal - EXECUTE COMMAND
.rdata:00012459
.rdata:0001245B C2 14 00        retn   14h ; DS:SI -> counted CR-terminated command string
.rdata:0001245B

```

Authors are trying to enumerate all useful ntoskrnl sections and for each of it call special function that performs searching *ZwAllocateVirtualMemory* by signatures on Windows 2000 or little harder on Windows XP.

```

.text:000119DD
.text:000119DD
.text:000119DD 0F B7 C3      jNextNtoskrnlSection:      ; CODE XREF: sub_11994+901j
.text:000119E0 6B C0 28      movzx  eax, bx
.text:000119E3 8D 14 30      imul  eax, 28h
.text:000119E6 E8 43 FD FF FF  lea    edx, [eax+esi]
.text:000119E6
.text:000119EB 84 C0          call   fnCheckSectionByFlagsame
.text:000119ED
.text:000119ED
.text:000119ED
.text:000119ED
.text:000119ED
.text:000119EF 8B 42 08      test   al, al
.text:000119F2 8B 4A 10      jz     short jNextIteration
.text:000119F5 3B C1          mov     eax, [edx+8]
.text:000119F7 73 02          mov     ecx, [edx+10h]
.text:000119F7
.text:000119F7
.text:000119F7
.text:000119F7
.text:000119F7
.text:000119F9 8B C8          cmp    eax, ecx
.text:000119F9
.text:000119F9
.text:000119FB
.text:000119FB
.text:000119FB 8B 42 0C      jnb   short loc_119FB ; CODE XREF: sub_11994+631j
.text:000119FE 03 45 F4      mov     eax, [edx+0Ch]
.text:00011A01 8D 55 FF      add     eax, [ebp+NTKernelBaseAddress]
.text:00011A04 52          lea    edx, [ebp-1]
.text:00011A05 8D 55 F8      push   edx
.text:00011A08 52          lea    edx, [ebp+var_8]
.text:00011A09 03 C8          push   edx
.text:00011A0B 51          add     ecx, eax
.text:00011A0C 50          push   ecx
.text:00011A0D 8D 45 D8      push   eax
.text:00011A10 50          lea    eax, [ebp+var_28]
.text:00011A11 C6 45 FF 01   push   eax
.text:00011A15 E8 6A FE FF FF  mov     [ebp+var_1], 1
.text:00011A15

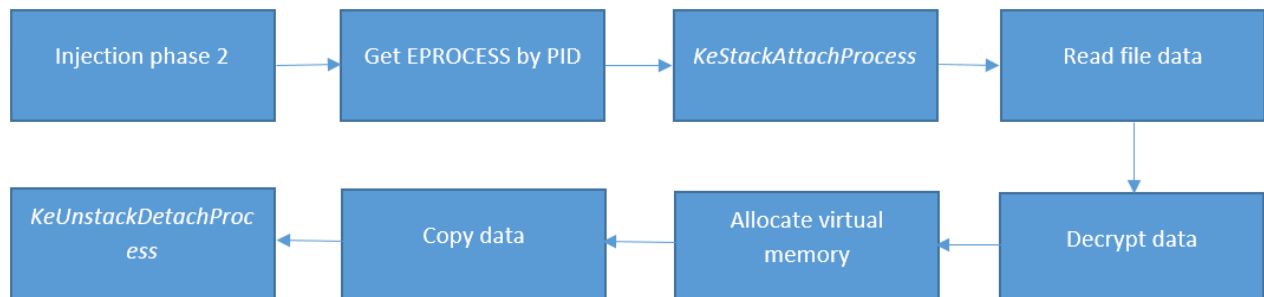
```

Main purpose of this Stuxnet rootkit is code injection. As we can see from its code, the driver tries to read configuration data of injection either from registry parameter *Data*, either from

file, if its name is present into malware sample. In analyzed sample, name of configuration file is absent. Injection configuration data is prepared by user mode part of malware. Injection mechanism was perfectly described by ESET in their paper. The driver performs injection into process in two phases: first phase is preparatory and second is major.

```
.text:000112B7 8D 75 D8          lea     esi, [ebp+var_28]
.text:000112BA 89 5D FC          mov     [ebp+var_4], ebx
.text:000112BD E8 7A 0A 00 00   call   fnLookupProcessAndAttachToIt
.text:000112BD
.text:000112C2 39 5D FC          cmp     [ebp+var_4], ebx
.text:000112C5 75 14            jnz     short loc_112DB
.text:000112C5
.text:000112C7 FF 75 14          push   [ebp+arg_C]
.text:000112CA 8B 4D 18          mov     ecx, [ebp+arg_10]
.text:000112CD FF 75 0C          push   [ebp+arg_4]
.text:000112D0 8B 55 10          mov     edx, [ebp+arg_8]
.text:000112D3 FF 75 08          push   [ebp+arg_0]
.text:000112D6 E8 95 FF FF FF   call   fnReadConfigFileAllocateMemAndWriteData_
.text:000112D6
.text:000112DB
.text:000112DB          loc_112DB:
.text:000112DB          ; CODE XREF: fnInjectionPhase2+1F1j
.text:000112DB          cmp     byte ptr [ebp+var_28], b1
.text:000112DE 74 0D            jz      short loc_112ED
.text:000112DE
.text:000112E0 8D 45 E4          lea     eax, [ebp+var_1C]
.text:000112E3 50              push   eax
.text:000112E4 8B 5D D8          mov     byte ptr [ebp+var_28], b1
.text:000112E7 FF 15 F0 23 01+  call   ds:KeUnstackDetachProcess
.text:000112E7 00
```

On second phase it tries to read content of file, decrypts it and injects it into process address space. File names for injection are stored into configuration file or registry parameter *Data*.



As we can see from the code analysis, authors have developed rootkit for injection malicious code into processes. Data for injection is prepared by Stuxnet user mode code. The driver registers handler for image load notify and performs injection into two phases. It also supports one IOCTL command for changing protection for virtual memory pages of process with help of *ZwProtectVirtualMemory*. For finding this unexported and undocumented function into ntoskrnl, it uses raw bytes search based on special signatures.

Driver 2

File name: Mrxnet.sys

SHA256: 0d8c2bcb575378f6a88d17b5f6ce70e794a264cdc8556c8e812f0b5f9c709198

File size: 17400 bytes

Signed: Yes

Timestamp: 2010-01-25 14:39:24

Device object name: none

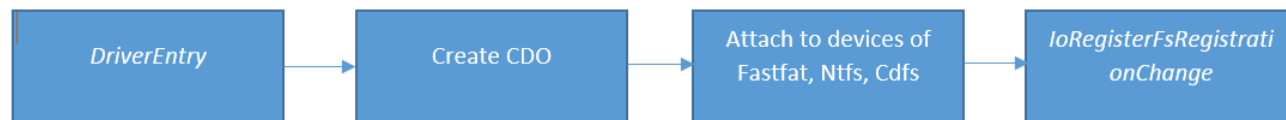
Main purpose: malicious files hiding

AV detection ratio: 51/61

Unlike first driver MRXCLS.sys, authors of Mrxnet.sys don't perform anti-analysis checks in the start function of driver. Mrxnet.sys is a FS filter driver that controls some file operations. The rootkit tries to hide some file types by controlling IRP_MJ_DIRECTORY_CONTROL request and removes information about it from buffer.

```
.text:00010765          jNextHandler:          ; CODE XREF: DriverEntry-188B↓j
.text:00010765 8B 0D 98 1F 01+      mov     ecx, DriverObject
.text:00010768 C7 04 08 86 04+      mov     dword ptr [eax+ecx], offset fnIrpMjHandlerCallNextDriver
.text:00010772 83 C0 04             add     eax, 4
.text:00010775 3D A4 00 00 00       cmp     eax, 0A4h
.text:0001077A 7E E9               jle     short jNextHandler
.text:0001077A
.text:0001077C A1 98 1F 01 00       mov     eax, DriverObject
.text:00010781 C7 40 6C 96 04+      mov     dword ptr [eax+6Ch], offset fnIrpMjHandlerFileSystemControl
.text:00010788 A1 98 1F 01 00       mov     eax, DriverObject
.text:0001078D C7 40 68 EC 04+      mov     dword ptr [eax+68h], offset fnIrpMjHandlerDirectoryControl
.text:00010794 E8 A9 FD FF FF       call    fnInitFastIoDispatch
.text:00010794
.text:00010799 E8 20 FF FF FF       call    fnAttachToExistingDevsOfFastfatNtfsCdfs
.text:00010799
.text:0001079E 68 EC 09 01 00       push   offset fnDriverFsNotification
.text:000107A3 FF 35 98 1F 01+      push   DriverObject
.text:000107A9 FF 15 08 1C 01+      call   ds:IoRegisterFsRegistrationChange
.text:000107AF 8B F0               mov     esi, eax
.text:000107B1 85 F6               test    esi, esi
.text:000107B3 7D 1D               jge     short loc_107D2
.text:000107B3
```

The rootkit initialization steps.



As we can see from code, the driver plays with two types of devices: firstly its own CDO (Control Device Object) that represents FS filter and secondly devices that were created to filter files related operations on specific volumes. In case of CDO, the rootkit dispatches widely known request IRP_MJ_FILE_SYSTEM_CONTROL and operation IRP_MN_MOUNT_VOLUME. This operation is used by Windows kernel in case of mounting new volume into a system. After got this request, the rootkit creates new device object, registers completion routine and attaches device to newly mounted device. This method allows for driver to monitor appearance in a system new volumes, for example, volume of removable drive.

```

.text:00010496          fnIrpMjHandlerFileSystemControl proc near
.text:00010496                                          ; DATA XREF: DriverEntry-1884↓j
.text:00010496  6A 08                push     8
.text:00010498  68 D8 1D 01 00      push     offset unk_11DD8
.text:0001049D  E8 22 13 00 00      call    fnSEH
.text:0001049D
.text:000104A2  83 65 FC 00         and     dword ptr [ebp-4], 0
.text:000104A6  8B 4D 0C            mov     ecx, [ebp+0Ch]
.text:000104A9  8B 41 60            mov     eax, [ecx+60h]
.text:000104AC  51                 push   ecx
.text:000104AD  FF 75 08            push   dword ptr [ebp+8]
.text:000104B0  80 78 01 01        cmp     byte ptr [eax+1], 1 ; IoStack->MinorFunction -- IRP_MN_MOUNT_VOLUME
.text:000104B4  75 0E              jnz     short jCallNextDriver
.text:000104B4
.text:000104B6  E8 4B 04 00 00      call    fnCreateDeviceAndAttachToDeviceStack
.text:000104B6
.text:000104BB
.text:000104BB          jRet_:
.text:000104BB          ; CODE XREF: fnIrpMjHandlerFileSystemControl+33↓j
.text:000104BB  C7 45 FC FE FF+    mov     dword ptr [ebp-4], 0FFFFFFEh
.text:000104C2  EB 1A              jmp     short jRet
.text:000104C2
.text:000104C4          ; -----
.text:000104C4
.text:000104C4          jCallNextDriver:
.text:000104C4          ; CODE XREF: fnIrpMjHandlerFileSystemControl+1E↑j
.text:000104C4  E8 7B 06 00 00      call    fnCallNextDriver
.text:000104C4
.text:000104C9  EB F0              jmp     short jRet_
.text:000104C9
.text:000104C9          fnIrpMjHandlerFileSystemControl endp

```

As you can see from the picture above, the driver also calls *IoRegisterFsRegistrationChange* I/O manager API for registering its handler that Windows kernel will call each time, when new file system driver CDO is registered into a system. In this handler, the driver creates new device and attaches it to passed CDO or removes device in case of file system driver deletion.

Major purpose of Mrxnet.sys driver is hiding Stuxnet malicious files. Windows kernel provides *ZwQueryDirectoryFile* API for requesting information about files in directory. This API function calls driver handler of IRP_MJ_DIRECTORY_CONTROL operation. So, the rootkit registers own IRP_MJ_DIRECTORY_CONTROL handler and sets completion routine when such request is passed through handler. In this completion routine it analyzes buffer with data and checks file names in it. It erases from buffer files with extension .LNK and .TMP. It also imposes additional restrictions on hiding. For example, in case of .LNK file, its size should be equal 0x104B.

```

.text:00011708 6A 04          push 4
.text:0001170A 8D 44 73 F8    lea  eax, [ebx+esi*2-8]
.text:0001170E 50            push  eax
.text:0001170F B8 98 1B 01 00  mov  eax, offset a_lnk ; ".LNK"
.text:00011714 E8 C1 FD FF FF  call  fnStrCmpi
.text:00011714
.text:00011719 84 C0          test  al, al
.text:0001171B 75 12          jnz   short jRemoveInfoFromBuf
.text:0001171B
.text:0001171D
.text:0001171D          loc_1171D:          ; CODE XREF: sub_11688+70↑j
.text:0001171D          ; sub_11688+79↑j ...
.text:0001171D FF 75 F4       push  [ebp+var_C]
.text:00011720 FF 75 F0       push  [ebp+var_10]
.text:00011723 56            push  esi
.text:00011724 8B F3         mov   esi, ebx
.text:00011726 E8 CB FE FF FF  call  fnCheckFileNameInList ; check on .TMP
.text:00011726
.text:0001172B 84 C0          test  al, al
.text:0001172D 74 1C          jz    short loc_1174B
.text:0001172D
.text:0001172F
.text:0001172F          jRemoveInfoFromBuf: ; CODE XREF: sub_11688+93↑j
.text:0001172F 8B 45 0C       mov   eax, [ebp+arg_4]
.text:00011732 85 C0          test  eax, eax
.text:00011734 74 2F          jz    short loc_11765
.text:00011734
.text:00011736 8B 4D FC       mov   ecx, [ebp+var_4]
.text:00011739 2B C8         sub   ecx, eax
.text:0001173B 51            push  ecx
.text:0001173C 03 C7         add   eax, edi
.text:0001173E 50            push  eax
.text:0001173F 57            push  edi
.text:00011740 FF 15 5C 1C 01+ call  ds:memmove
.text:00011746 83 C4 0C       add   esp, 0Ch

```

check .LNK extension

check .TMP extension

erase info

It should be noted that such technique of files hiding were described in famous book "Rootkits: Subverting the Windows kernel" by Hoglund, Butler.

Driver 3

File name: Jmidebs.sys

SHA256: 63e6b8136058d7a06dfff4034b4ab17a261cdf398e63868a601f77ddd1b32802

File size: 25552 bytes

Signed: Yes

Timestamp: 2010-07-14 09:05:36

Device object name: \Device\{3093983-109232-29291}

Main purpose: code injection

AV detection ratio: 50/61

This driver is pretty similar to MRxCls.sys and serves only for code injection into processes. The following properties distinguish it from original MRxCls.sys.

- New device object name - \Device\{3093983-109232-29291}
- New registry service name - jmidebs
- New registry service parameter name (injection data) - IDE
- New IOCTL code for reading (caching) configuration data
- New constants in decryption routine.

Decrypted strings.

\REGISTRY\MACHINE\SYSTEM\CurrentControlSet\Services\jmidrebs
IDE
\Device\{3093983-109232-29291}

```
; ecx->data; esi->size  
; Attributes: bp-based frame  
  
fnDecryptData proc near ; CODE XREF: DriverEntry+7E1p  
; sub_10070+731p ...  
  
var_10 = dword ptr -10h  
var_C = dword ptr -0Ch  
var_8 = dword ptr -8  
var_4 = dword ptr -4  
  
push ebp  
mov ebp, esp  
sub esp, 10h  
mov ecx, eax  
xor edx, 0D4114896h  
xor eax, 0A36ECD00h  
mov [ebp+var_4], esi  
shr [ebp+var_4], 1  
push ebx  
mov [ebp+var_10], edx  
mov [ebp+var_C], eax  
mov [ebp+var_8], 4  
push edi
```

MRxCls.sys

```
; ecx->data; esi->size  
; Attributes: bp-based frame  
  
fnDecryptData proc near ; CODE XREF: fnInitDriverEntry+1c  
; sub_10C22+881p ...  
  
var_10 = dword ptr -10h  
var_C = dword ptr -0Ch  
var_8 = dword ptr -8  
var_4 = dword ptr -4  
  
push ebp  
mov ebp, esp  
sub esp, 10h  
mov ecx, eax  
xor edx, 0BADF00Dh  
xor eax, 0EADBEEFh  
mov [ebp+var_4], esi  
shr [ebp+var_4], 1  
push ebx  
mov [ebp+var_10], edx  
mov [ebp+var_C], eax  
mov [ebp+var_8], 7  
push edi
```

Jmidrebs.sys

The rootkit contains additional IOCTL function, which specializes in caching configuration data. This data are used for injection malicious Stuxnet code.

```
.text:00010B68 fnDispatchDeviceControlRequest proc near  
.text:00010B68 ; CODE XREF: fnDispatchIrpMjFunctions+301p  
.text:00010B68  
.text:00010B68 var_4 = dword ptr -4  
.text:00010B68 55 push ebp  
.text:00010B69 8B EC mov ebp, esp  
.text:00010B6B 51 push ecx  
.text:00010B6C 8B 41 60 mov eax, [ecx+60h]  
.text:00010B6F 8B 40 0C mov eax, [eax+0Ch]  
.text:00010B72 2D 00 38 22 00 sub eax, 223800h  
.text:00010B77 74 20 jz short loc_10B99  
.text:00010B79 83 E8 04 sub eax, 4  
.text:00010B7C 74 07 jz short jDispatchNewIOCTLFunc  
.text:00010B7C  
.text:00010B7E B8 02 00 00 C0 mov eax, 0C0000002h  
.text:00010B83 C9 leave  
.text:00010B84 C3 retn  
.text:00010B84  
.text:00010B85 ; -----  
.text:00010B85  
.text:00010B85 jDispatchNewIOCTLFunc: ; CODE XREF: fnDispatchDeviceControlRequest+141j  
.text:00010B85 56 push esi  
.text:00010B86 8D 75 FC lea esi, [ebp+var_4]  
.text:00010B89 E8 54 FF FF FF call fnInit  
.text:00010B89  
.text:00010B8E 8B 45 FC mov eax, [ebp+var_4]  
.text:00010B91 E8 8C 00 00 00 call fnReadConfigurationData  
.text:00010B91
```

new IOCTL function

```

.text:00010A57 8D 75 FC          lea     esi, [ebp+var_4]
.text:00010A5A E8 83 00 00 00   call   fnInit
.text:00010A5A                                mov     eax, [ebp+var_4]
.text:00010A5F 8B 45 FC          call   fnReadConfigurationData
.text:00010A62 E8 BB 01 00 00   test   eax, eax
.text:00010A67 85 C0            jnz    short loc_10A86
.text:00010A69 75 1B            lea     esi, [ebp+var_8]
.text:00010A6B 8D 75 F8          call   sub_10B36
.text:00010A6E E8 C3 00 00 00   lea     esi, [ebp+var_8]
.text:00010A6E                                call   sub_1056C
.text:00010A73 8D 75 F8          push   offset fnLoadImageNotify
.text:00010A76 E8 F1 FA FF FF   call   ds:PsSetLoadImageNotifyRoutine
.text:00010A76                                ; CODE XREF: fnPrepareForCodeInjection+1B1j
.text:00010A7B 68 80 0F 01 00   ; fnPrepareForCodeInjection+2F1j
.text:00010A80 FF 15 AC 25 01+
.text:00010A80 00
.text:00010A86
.text:00010A86     loc_10A86:
.text:00010A86     ; CODE XREF: fnPrepareForCodeInjection+1B1j
.text:00010A86     ; fnPrepareForCodeInjection+2F1j

```

first config data caching

for second client uses IOCTL 0x223804

Driver 4

File name: MRxCls.sys

SHA256: 1635ec04f069ccc8331d01fdf31132a4bc8f6fd3830ac94739df95ee093c555c

File size: 26616 bytes

Signed: Yes

Timestamp: 2009-01-01 18:53:25

Device object name: \Device\MRxClsDvX

Main purpose: code injection

AV detection ratio: 50/61

This sample is identical to driver 1, but signed with digital certificate. Both samples have identical timestamp value in PE header and identical code inside.

Conclusion

As we can see from the analysis, authors of Stuxnet Ring 0 part were interested in code injection and malicious files hiding. Driver MRxCls.sys has two instances, one unsigned and another with digital signature. Both drivers are identical and contain same compilation date. Driver Jmidebs.sys was compiled later than these two and I can call it "MRxCls.sys v2", because it contains some differences inside, but serves for same purpose. Driver Mrxnet.sys is a typical legacy FS filter driver that is used by attackers for hiding files in Windows.

| | Driver 1 | Driver 2 | Driver 3 | Driver 4 |
|-----------------|----------------------------|----------------------------|------------------------|------------------------|
| File name | MRxCls.sys | MRxCls.sys | Mrxnet.sys | Jmidebs.sys |
| File size | 19840 bytes | 26616 bytes | 17400 bytes | 25552 bytes |
| Signed | No | Yes | Yes | Yes |
| Timestamp | 2009-01-01 18:53:25 | 2009-01-01 18:53:25 | 2010-01-25 14:39:24 | 2010-07-14 09:05:36 |
| Main purpose | Code injection | Code injection | Malicious files hiding | Code injection |
| Uses encryption | Yes | Yes | No | Yes |
| Device name | MRxClsDvX | MRxClsDvX | CDO/Unnamed | {3093983-109232-29291} |