

Unpacking Smokeloder and Reconstructing PE Programatically using LIEF

 m.alvar.es/2020/06/unpacking-smokeloder-and.html

This article holds notes on my experience unpacking a *Smokeloder* 2020 sample. The unpacked payload is further used for composing a valid *PE* file. The outcome is a *PE32* executable containing clean code ready for reversing.

First things first, here is the sample used in this research:

Size	308.17 KB (315568 bytes)
Type	PE32 executable for MS Windows (GUI) Intel 80386 32-bit
First seen	2020-03-06 21:45:11
md5	c067e0a2d7fc6092bb77abc7f7156b60
sha1	52f68073caec0fd424c7cbaaed5f5221d7103d20
sha256	25959cfe4619126ab554d3111b875218f1dbfadd79eed1ed0f6a8c1900fa36e0

You can find it in VirusTotal [\[1\]](#).

This sample does regular already documented *Smokeloder* checks before unpacking the main payload, such as:

- checks if the process is running in the context of a debugger using "*kernel32.isDebuggerPresent*" function [\[2\]](#);
- makes a copy of *ntdll.dll*, loads it and uses it instead. This technique helps to evade some sandboxes and has been described already in this article here [\[3\]](#);
- looks for specific patterns in registry keys to check if the sample is running under a virtualised environment.

It also performs a small profiling of the hosting machine in order to decide which payload to inject. *Smokeloder* has specific code for both main architectures x86 and x64. In this article, we gonna unpack the x86 payload of the above mentioned sample.

Smokeloder has been using various techniques to inject its final payload into the user file management process "*explorer.exe*". The sample analysed uses *RtlCreateUserThread* approach in order to copy the final payload to the targeted process. This injection method is better described in this Endgame/Elastic article [\[4\]](#).

So our game plan is:

1. pause execution before the unpacked payload is executed by "*explorer.exe*";
2. transplant this code to a dummy PE shell;
3. fix PE header values and section boundaries;
4. patching *Smokeloder* code preamble;
5. test unpacked *Smokeloder* PE;
6. how to do all this programatically using LIEF [\[5\]](#).

1.0 - FETCHING MAIN PAYLOAD

Smokeloder 1st stage decompresses its payload using *ntdll.RtlDecompressBuffer* [\[6\]](#) after few anti-analysis checks described above. It does not call this function from the initially loaded *ntdll.dll* but from a copy of it loaded afterwards. So breakpoints should be set after the binary loads the copy of *ntdll.dll*. Figure 01 presents a screenshot of this specific code *IDA*.

```

seg001:004013E2    lea     eax, [ebp+var_2]
seg001:004013E5    mov     [ebp+buffer_size], edx
seg001:004013E8    lea     edx, [ebp+buffer_size]
seg001:004013EA    push   4                ; Protect
seg001:004013EA    push   3000h           ; AllocationType
seg001:004013EF    push   edx             ; RegionSize
seg001:004013F0    push   ecx             ; ZeroBits
seg001:004013F1    push   eax             ; BaseAddress
seg001:004013F2    push   0FFFFFFFFh     ; ProcessHandler
seg001:004013F4    call   ds:ZwAllocateVirtualMemory
seg001:004013FA    test   eax, eax
seg001:004013FC    jnz    short loc_411423
seg001:004013FE    mov     eax, [ebp+buffer_addr]
seg001:00401401    lea     ecx, [ebp+var_4]
seg001:00401404    lea     edx, [ebp+buffer_size]
seg001:00401407    push   ecx             ; FinalUncompressedSize
seg001:00401408    push   edi             ; CompressedBufferSize
seg001:00401409    push   esi             ; CompressedBuffer
seg001:0040140A    push   dword ptr [edx]; UncompressedBufferSize
seg001:0040140C    push   eax             ; UncompressedBuffer
seg001:0040140D    push   2                ; COMPRESSION_FORMAT_LZNT1
seg001:0040140F    call   ds:RtlDecompressBuffer
seg001:00401415    test   eax, eax
seg001:00401417    jnz    short loc_411423
seg001:00401419    mov     eax, [ebp+var_4]
seg001:0040141C    mov     ecx, [ebp+arg_C]
seg001:0040141F    mov     [ecx], eax
seg001:00401421    jmp    short loc_41142A

```

Figure 01: *Smokeloder* first stage decompression code

This code allocates a buffer with *0x2D000* bytes using *ZwAllocateVirtualMemory* which stores the main decompressed payload [7]. This code is still transformed before being injected into "*explorer.exe*". The following steps are performed during injection:

- fetches *explorer.exe* PID by calling *GetShellWindow* and *GetWindowThreadProcessId*;
- sections and maps are created in the current and remote processes using *ZwCreateSection* and *ZwMapViewOfSection*;
- main payload is copied to local section and reflected in the remote section;
- data section is created in the remote process for holding parameters and dynamically created Import Table;
- A new thread is created in the remote process by invoking *RtlCreateUserThread*.

So, at this point, you could ask me: what is the relevance of describing all these call names to the final goal of this article? the answer is: so you can reproduce exactly what I'm describing in here. :D

Next step is setting up a break point in *RtlCreateUserThread* (from the copy of *ntdll.dll*) and dump the final payload. It is also necessary to take note of few important addresses: (i) entry point of the thread created in the remote processes and (ii) base addresses for injected code in virtual process.

Figure 02 shows a screenshot of *IDApro* showing the call to *RtlCreateUserThread* (where we should pause the execution).

```

seg001:004017AC    loc_4017AC:                ; CODE XREF: sub_401468
seg001:004017AC    pop     esi
seg001:004017AD    mov     ecx, [esi+28h]
seg001:004017B0    add     ecx, [ebp+var_40]
seg001:004017B3    xor     edi, edi
seg001:004017B5    lea     eax, [ebp+var_4]
seg001:004017B8    mov     [eax], edi
seg001:004017BA    push   edi             ; ClientID
seg001:004017BB    push   eax             ; ThreadHandle
seg001:004017BC    push   [ebp+var_3C]    ; StartParameter
seg001:004017BF    push   ecx             ; StartAddress
seg001:004017C0    push   edi             ; StackCommit
seg001:004017C1    push   edi             ; StackReserved
seg001:004017C2    push   edi             ; StackZeroBits
seg001:004017C3    push   edi             ; CreateSuspended
seg001:004017C4    push   edi             ; SecurityDescriptor
seg001:004017C5    push   [ebp+var_C]    ; ProcessHandle
seg001:004017C8    call   dword ptr ds:RtlCreateUserThread
seg001:004017D7

```

Figure 02: Call to *RtlCreateUserThread* after injecting code into remote process

By stopping the execution on this call we can collect all data we need to move on to the next step:

Base address code	0x02060000
Base address data	0x00B60000

Data payload	<code>a01751fb6eb3f19d9b010818bbecc23c</code> [8]
Code payload	<code>2547231b4ae82ea9e395fb0c8a308982</code> [9]
Code entry point	<code>0x02061734</code>

Code payload is the final unpacked *Smokeloader* code adjusted to run on Virtual Address with base equal to `0x02060000`. The created thread receives the base address of the data segment (`0x00B60000`) as parameter ("*StartParameter*" parameter of "*RtlCreateUserThread*" call).

2.0 - TRANSPLANTING PAYLOAD TO PE

Smokeloader loads all resources necessary to its execution dynamically. This article [here](#) [10] describes how *Smokeloader* builds up its import table and how to prepare patch an IDB to overcome this technique before starting reversing. So this main payload does not need any specific setup of *imports*.

In this section we will use 010 Hex Editor [11] to transplant a PE header from a random executable. 010 Hex Editor has a PE format template [12]. Although any other valid PE32 binary could be used in this experiment, we used a PE header extracted from an executable listed in this Sotirov's blog post [13][14].

Smokeloader code payload has `0x1000` null bytes at offset zero, so we copied the first `0x1000` bytes containing the PE header from *tinype.exe* to this region.

Coincidentally, *.text* section will be already pointing to the beginning of the our payload at offset `0x1000`. Probably the malware author just wiped out the PE header before creating the payload and left the null bytes there. Next step is to paste all `20480` bytes (`0x5000`) of our data payload in the offset `0x4400`.

Figure 03 shows the new layout of our binary containing the PE header in the beginning followed by `0x3400` bytes of code payload (at offset `0x1000`) and finally `0x5000` bytes of data from our data payload (at offset `0x4400`).

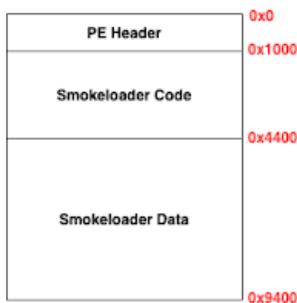


Figure 03: Initial layout of new handcrafted PE binary

3.0 - FIXING PE HEADERS

It is time to adjust our implanted PE header manually using *010 Hex Editor*. At this point, all fields in this header are still set up according to "*tinype.exe*". From now on, we gonna use this schematic as reference to PE header internal structures [15].

The first adjustment is to change the number of sections to 2 for holding code and data. This field is located in the "*COFFHeader.NumberOfSections*". Now our binary will list only 2 sections named *.text* and *.rdata* we can rename this second one to *.data* by changing "*SectionHeaders[1].Name*".

Next step is make sure that both sections have correct permissions. "*SectionHeaders[0].Characteristics*" (*.text*) should have *CODE*, *EXECUTE* and *READ* flags active and "*SectionHeaders[1].Characteristics*" (*.data*) should have the

INITIALIZED_DATA, *READ* and *WRITE* flags active. Still on *SectionHeaders*, we can setup the bounds and virtual addresses. "*SectionHeaders[0].SizeOfRawData*" should be set to *0x3400* (13312 Bytes), "*SectionHeaders[0].PointerToRawData*" should be set to *0x1000* and finally "*SectionHeaders[0].VirtualAddress*" should be set to *0x1000*. For "*SectionHeaders[1]*" ("*.data*") we gonna set "*SizeOfRawData*" to *0x5000*, "*PointerToRawData*" to *0x4400* and "*VirtualAddress*" to *0x5000*. These changes means that these sections will be mapped in memory in *base_address* (defined in the *OptionalHeader*) shifted by each section Virtual Addresses offset. There is an Union inside these section headers called "*PhysicalAddress*" and "*VirtualSize*", these fields should hold the same value as "*SizeOfRawData*".

Figure 04 shows a diagram of a Section header. Each section in the binary has an instance of this header associated to it.

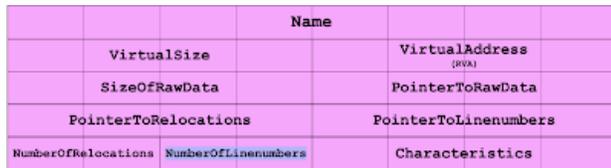


Figure 04: PE Section Header

Now we need to adjust few fields in the *Optional Header*. In this header we will need to change the following fields:

ImageBase	Virtual Address where binary will be mapped	0x02060000
SizeOfCode	size of .text section	0x3400 bytes
SizeOfInitializedData	size of .text and .data sections together	0x8400 bytes
AddressOfEntryPoint	offset of the entry point code	0x1734
BaseOfCode	.text section Virtual Address	0x1000
BaseOfData	.data section Virtual Address	0x5000
SectionAlignment	Virtual Addresses have to be multiple of this value	0x1000
FileAlignment	file offsets have to be multiple of this value	0x200
SizeOfImage	total size of binary headers + sections	0x9400
Checksum	PE file checksum - use PE Explorer [16] or Hiew [17] to calculate this value	-----

"*ImageBase*" has to match the base of the code section we dumped from "*explorer.exe*" (*0x02060000*). As we will not export or import anything all "*Data Directories*" inside the *Optional Header* can be *zeroed* as well.

Summarising the whole process:

1. Transplanting PE header from a dummy PE;
2. Fix sections sizes, boundaries, permissions and Virtual Addresses in *SectionHeaders*;
3. Setup section contents;
4. Setup *Optional Header* fields;
5. Setup PE checksum;

Here is the version of our binary after following up all steps described above [18]. This binary is a valid executable and we can load it in any debugger or disassembly but we still need to change one last thing before call it a valid unpacked *Smokeloader* sample.

4.0 - PATCHING BINARY

Figure 05 shows our reconstructed PE loaded in *IDApro* paused on the correct Entry Point.

```
.text:02061734      public start
.text:02061734      start          proc near
.text:02061734      arg_0          = dword ptr 8
.text:02061734      push         ebp
.text:02061735      mov         ebp, esp
.text:02061737      mov         ecx, [ebp+arg_0] ; <== Copying thread parameter to ECX
.text:0206173A      call        sub_2061743
.text:0206173F      pop         ebp
.text:02061740      retn        4
.text:02061740      start          endp
.text:02061743      ; ===== SUBROUTINE =====
.text:02061743      sub_2061743    proc near          ; CODE XREF: start+61p
.text:02061743      push        esi
.text:02061743      mov         esi, ecx
```

Figure 05 - Reconstructed PE paused on Entry Point

We can notice that the entry point function receives an argument (*0x02061737*) and loads it into *ECX* and then calls another function located in *0x02061743* which is just below the current function. This argument is the address of the data segment. This data segment will be used for various tasks during *Smokeloader* execution including holding the dynamically created import table.

If we execute this file without a valid value in *ECX* it will break when the main payload tries to write into the data segment (invalid address in *ECX*). Figure 06 shows what happens when we try to execute our binary the way it is right now.

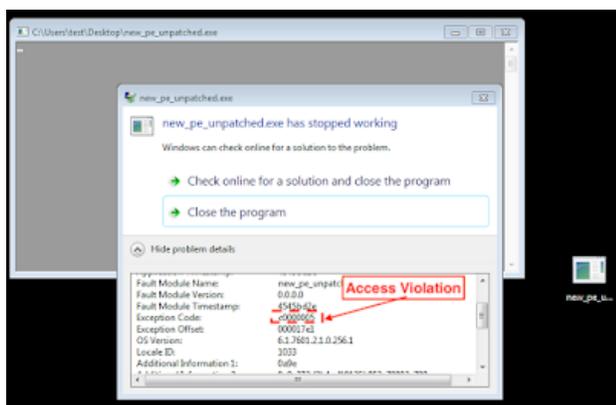


Figure 06 - Access Violation exception when executing unpatched reconstructed binary

The plan now is to patch this binary to load the correct address of the data segment into *ECX* before calling "*sub_2061743*". Since both functions are consecutive and the function on top does not do much - we gonna replace all 15 bytes of this function (*0x02061743 - 0x02061734*). Figure 07 shows the new patched code.

```

.text:02061734                                     public start
.text:02061734                                     start      proc near
.text:02061734 68 00 50 06 02                               push     offset asc_2065000 ; ""
.text:02061739 59                               pop      ecx
.text:0206173A 90                               nop
.text:0206173B 90                               nop
.text:0206173C 90                               nop
.text:0206173D 90                               nop
.text:0206173E 90                               nop
.text:0206173F 90                               nop
.text:02061740 90                               nop
.text:02061741 90                               nop
.text:02061742 90                               nop
.text:02061743 56                               push     esi
.text:02061744 8B F1                             mov     esi, ecx

```

Figure 07 - Code after patching

In this new code the entry point remains the same. We can see that we loaded *ECX* with the address of the data segment by using the push and pop instructions and then we filled the rest of the remaining bytes with *NOP (0x90)*. We can see the beginning of the second function at the same address as before (*0x02061743*). Of course there are many ways to achieve this same result but this was the simplest approach we could think of.

The final step is to update the PE checksum field inside the *Optional Header* again and we will have a fully unpacked *Smokeloader* sample. Here are the last version of our reconstructed binary [19]:

File name	new_pe_patched.bin
File type	PE32 executable for MS Windows (console) Intel 80386 32-bit
Size	37888 bytes
md5	f401109ae24aaf47dce75266ffc049f8
sha1	49e7ed68b9569e0e987da71b3c678974d8ed7c81
sha256	cd42f017913034d527d90a84feebcde015e714baa03714c83f80608555e52386

5.0 - TESTING RECONSTRUCTED PE

For testing our branding new reconstructed PE we ran it into Cuckoo sandbox [20] to analyse its behaviour [21]. As we can see in figure 08 and 09, the binary was executed properly and we got it checking in and contacting its controllers.

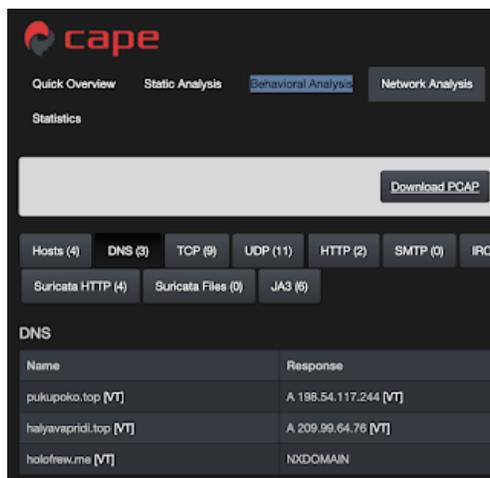


Figure 08 - Reconstructed sample connecting back to Controllers

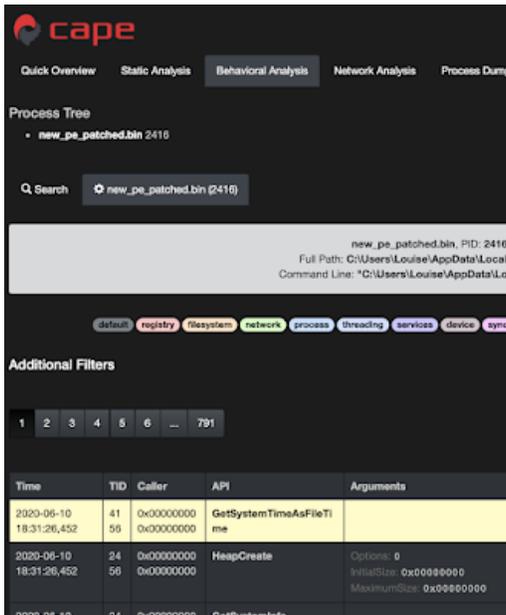


Figure 09 - List of API calls intercepted by Cuckoo

As we can see we got the sample connecting back to three controller *URLs* and many pages of intercepted *API* calls in the behavioural analysis. This is an indication that our unpacked and reconstructed *Smokeloader* sample is functional.

6.0 - DOING IT ALL PROGRAMATICALLY

So far we described how to unpack *Smokeloader* main payload and how to manually reconstruct a valid PE file out of this. Now we will automate what we did manually by transplanting a PE header from a dummy binary ("*tinype.exe*"). The Python library used in this experiment is *LIEF* [5]. We extended this example called "Create a PE from scratch" they have in their official documentation [22].

The following code does exactly the same as the manual approach but using *LIEF*.

The final binary generated by *LIEF* is:

File name	unpacked_smokeloader.exe
File type	PE32 executable for MS Windows (console) Intel 80386 32-bit
Size	34816 bytes
md5	a0aebc61bc89208be0585eca4d1ed00c
sha1	ea2f3c914dec6bb36832abc313b3fce826cdec0
sha256	0247de510507792fcbf425fab9dbbc2f067c25dc7e4e80a958d1ebfb0505f6e6

We uploaded it for testing to Virustotal [23] and CAPE sandbox [24] and is a valid unpacked *Smokeloader* PE32 executable.

REFERENCES:

- [1] <https://www.virustotal.com/gui/file/25959cfe4619126ab554d3111b875218f1dbfadd79eed1ed0f6a8c1900fa36e0/details>
- [2] <https://docs.microsoft.com/en-us/windows/win32/api/debugapi/nf-debugapi-isdebuggerpresent>
- [3] <https://malwareandstuff.com/examining-smokeloaders-anti-hooking-technique/>
- [4] <https://www.elastic.co/blog/ten-process-injection-techniques-technical-survey-common-and-trending-process>
- [5] <https://github.com/lief-project/LIEF>
- [6] <https://docs.microsoft.com/en-us/windows-hardware/drivers/ddi/ntifs/nf-ntifs-rtldecompressbuffer>
- [7] <https://www.virustotal.com/gui/file/756dd799c8195b98f295baa210ac1807f7d1d86de2736f559c76ce1c7816d0ee/detection>
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- [10] <http://security.neurolabs.club/2019/10/dynamic-imports-and-working-around.html>
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- [12] https://www.sweetscape.com/010editor/repository/templates/file_info.php?file=EXE.bt
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- [16] <http://www.heaventools.com/>
- [17] <http://www.hiew.ru/>
- [18] <https://www.virustotal.com/gui/file/3e474e495e11716f6eab40ee3c353602da2683dbefd900f6d90dfcedff2fa93d>
- [19] <https://www.virustotal.com/gui/file/cd42f017913034d527d90a84feebcde015e714baa03714c83f80608555e52386/detection>
- [20] <https://www.capesandbox.com/>
- [21] <https://www.capesandbox.com/analysis/7521/#behavior>
- [22] https://lief.quarkslab.com/doc/latest/tutorials/02_pe_from_scratch.html
- [23] <https://www.virustotal.com/gui/file/0247de510507792fcbf425fab9dbbc2f067c25dc7e4e80a958d1ebfb0505f6e6>
- [24] <https://www.capesandbox.com/analysis/7643/>